

A New Joining Algorithm for Multi-hop Non-mobile Wireless Sensor Networks

M. Çibuk¹ , D. Arı² , F. Ağgün³

¹Bitlis Eren University, Department of Computer Engineering, Bitlis, Turkey, (mcibuk@beu.edu.tr)

²Bitlis Eren University, Department of Computer Engineering, Bitlis, Turkey, (dari@beu.edu.tr)

³Bitlis Eren University, Department of Computer Technologies, Bitlis, Turkey, (faggun@beu.edu.tr)

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Corresponding author: M. Çibuk

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ABSTRACT

Wireless Sensor Networks (WSN) are divided into two categories as single-hop and multi-hop according to the connection states of the sensor nodes. In the single-hop WSNs, the sensor nodes communicate directly with the management unit (coordinator node). On the other hand, in multi-hop WSNs, the sensor nodes, which are out of the coverage area of the coordinator node communicate with the coordinator over other joined sensor nodes. Multi-hop WSNs preferred for complex applications where there are many sensor nodes. In such applications, connecting to the network and maintaining the continuity in the network is difficult and requires very complicated algorithms for sensor nodes.

In this study, a new network-joining algorithm for multi-hop WSNs has been proposed and designed. According to the algorithm, there is a management unit which is named as coordinator node (CN) for maintaining the network connectivity and the nodes which are the members of the network sends network joining requests over a common contention-based channel for joining the network. If there is a coordinator node in the coverage area, the sensor node joins directly to the network and coordinator allocates an appropriate channel to the new node. If the channel has multiple users, it is used as time-shared. If the sensor node is out of coverage area, directly joining to the network is impossible. In this case, the sensor node sends a relay request to the other joined sensor nodes. The available sensor node responds to this relay request via handshake. The sensor node uses the relay node's channels as shared by using time-sharing methods. Therefore, the new node joins the network from the relay node's channel. The main goal of this algorithm is to speed up joining of sensor nodes to the networks in multi-hop WSNs. Also, another goal is maximizing the continuity of the sensor nodes in the network. Thus, the algorithm improves performance and ensures data transfer continuity of WSNs.

1. INTRODUCTION

Wireless Sensor Networks (WSNs) consists of Sensor Nodes (SNs) connected wirelessly with each other. In the WSNs, there are sensor nodes and a coordinator node as basic network elements. The Coordinator Node (CN) organizes the other network elements. The most basic feature of sensor nodes is to sense environmental data, to process sensed data and to transmit processed data[1]–[3]. The required sensor node count changes from application to application. For example, in order to track environmental imaging and natural catastrophes, because of the land structure, requires a large number of sensor nodes. It is obvious that increasing the number of sensor nodes in the network makes more complicated the network. Complex structured WSNs use a multi-hop network structure as a connection method. In such networks, as shown in Figure 1, the sensor nodes extend coverage of the network by communicating with each other [4], [5]. In multi-hop networks, the sensor nodes perform not

only their own sensing tasks but also for its neighbor nodes the relay function. Multi-hop WSNs have advantages such as coverage, high data transmission rate, low cost[6]. Thanks to these networks, which spread over large areas, data can be collected by consuming less energy. Thus, the lifetime of the network extends[7].

On the contrary of the above mention advantages of the multi-hop WSN, there are some difficulties. Difficulties which foreseen in the multi-hop WSN design are negative factors such as congestion[1], end-to-end delay, hidden node, simultaneous communication, a disconnection of nodes from the network, topology changes[4], and fairness[6]. These difficulties have a negative impact on the performance of the network. When looking at the literature[8]–[11], the existence of studies can be observed on which contains the solution of these difficulties. While the studies, which do for the post-setup phase of the network are more frequent, the studies for network initialization phase seems to be less.

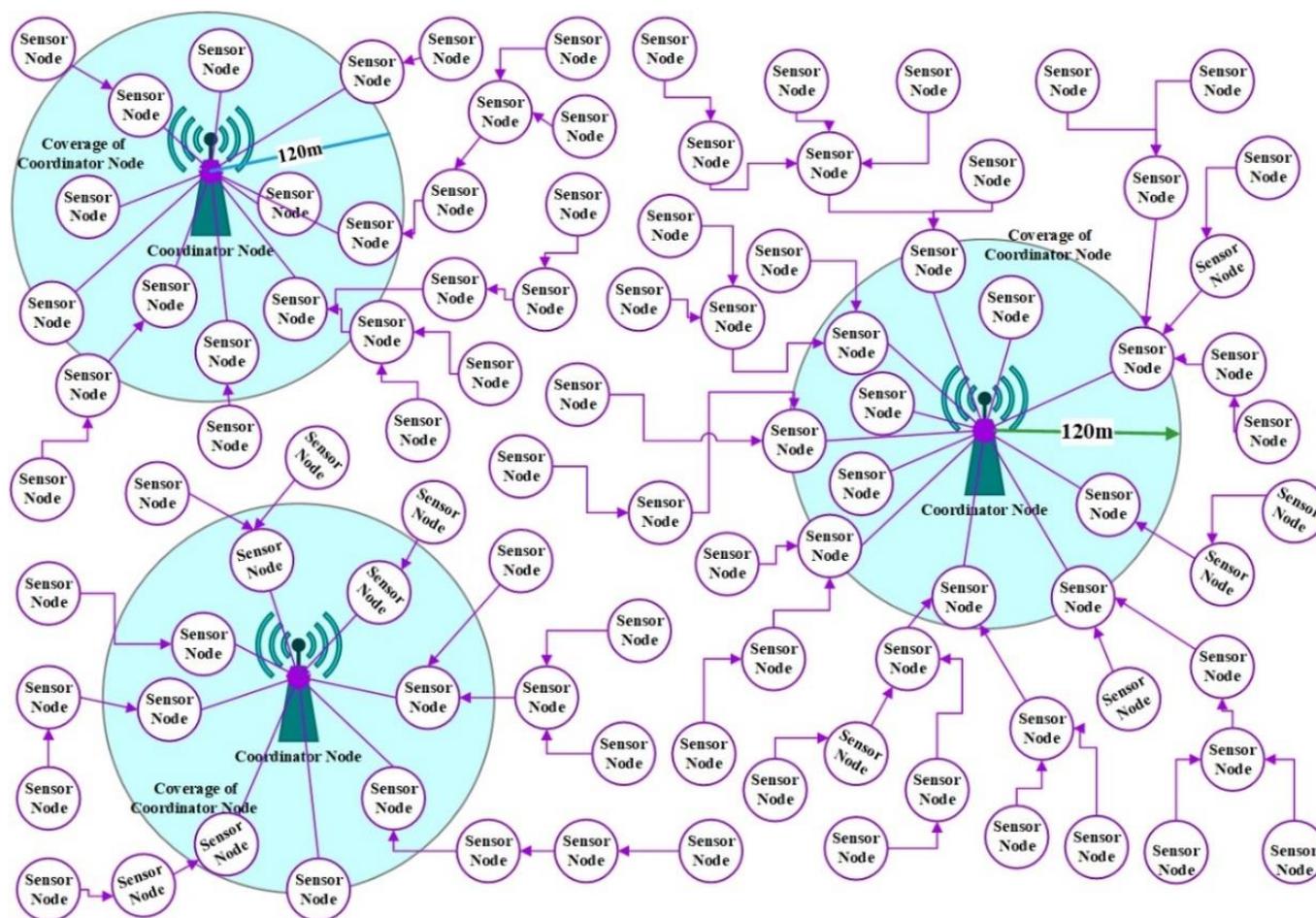


Figure 1. Connection state of Multi-hop Wireless Sensor Networks

In this study, a new joining network algorithm have been developed for the multi-hop non-mobile WSNs. According to the algorithm, the sensor nodes have sensing functions as well as a relay function that for other nodes can be used when needed. Furthermore, in this algorithm, a multi-channel (FDMA) structure has been used to minimize the collision. The nodes use a common contention-based (CSMA-CA) channel for their request to join the network, after which for the communication the nodes uses the channel which allocated by the CN. The sensor node directly joins the network if it is in the coverage area of the CN. If the sensor node is not in the coverage area of CN, the nodes send a request in order to join the network to the nearest sensor node, which has joined before. Neighbor sensor node, which received this network-joining request, sends a request to the CN. If CN accepts to request, the node sets the request-sender node as its relay node and use channels of the relay node. The nodes communicate with the CN as time-shared (TDMA) by using relay node channels. The performance of this algorithm has been tested based on node counts at random form of nodes in the Riverbed Modeler simulation environment.

2. A NEW JOINING ALGORITHM FOR MULTI-HOP NON-MOBILE WIRELESS SENSOR NETWORKS

According to the algorithm, firstly, the nodes listen to the *CH_REQUEST* channel and attempts to obtain the *CH_REQUEST* channel by using the CSMA-CA method, as

shown in Figure 2. *CH_REQUEST* is a common channel, which nodes use to send *CTRL_PKT* (control packet, network-join request) packets to the CN. When the node obtained the channel, it sends a *CTRL_PKT* to the CN. Later, the node starts to listen to the *CH_SCHEDULE* channel that CN uses to send the *SCH_PKT*s which tell conditions of nodes in the network to the SNs. After this stage, the node starts a counter to hold the *Count* for not scheduled and a *wait_time* counter for determining the time without getting any packet. If there is *SCH_PKT* on the channel and a channel allocation is present for itself, the node will be connected to the CN directly. Otherwise, if there is *SCH_PKT* on the channel but channel allocation is not for itself counter will be incremented and it will be returned to listen to *CH_SCHEDULE* channel state until the counter exceeds the *threshold_count* threshold. If there is not any packet on the *CH_SCHEDULE* channel within the *wait_time* the node sets *CTRL_PKT*s relay/request bit as 1 and starts to broadcast a this packet to the neighboring sensor nodes.

The neighboring node that received this packet will send a *RRQUEST_PKT* (relay request packets which supply to join as a relay to the network) to the CN. The relay node becomes the parent of the sensor node, which wants to join the network. After this stage, the relay node starts to send the *SCH_PKT*s which coming from CN to all child nodes by using the *CH_SCHEDULE* channel. The searching relay node or connecting CN directly period continues as described below. After this step, the node repeats the same phase as the phase of direct CN connection that mentions above, as shown in Figure 2.

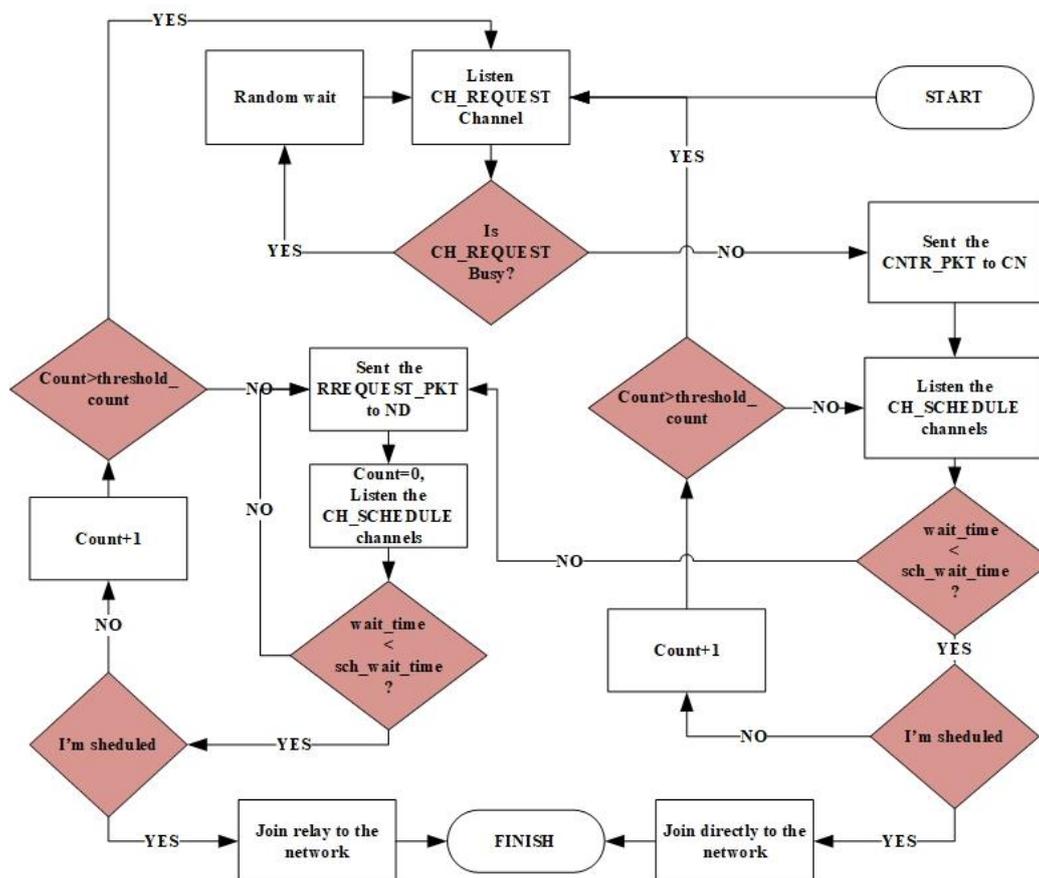


Figure 2. Flowchart of the algorithm

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Start
Initial_Values_Set, Count=0, Time_out,
wait_time
RREQUEST_PKT, SCH_PKT, CTRL_PKT,
CH_SCHEDULE, CH_REQUEST, CN, ND;
Step:1 Listen CH_REQUEST channel
Step:2 If (Is CH_REQUEST channel busy?)
    Random_wait;
    GoTo Step1;
    Else
        Sent CTRL_PKT to CN
        GoTo Step4;
Step:3 While(Listen CH_SCHEDULE channel){
    If (Is there SCH_PKT)
        GoTo Step4;
    Else
        GoTo Step5
Step:4
    If(I'm sheduled)
        GoTo Step7;
    Else
        Count++;
        If(Count=> threshold_count)
            GoTo Step5;
        Else
            GoTo Step1;
}
Step:5 sets CTRL_PKT's relay/request bit TRUE
Send CTRL_PKT
Count=0;
While (Listen CH_SCHEDULE channel){
    If (Is there SCH_PKT)
        GoTo Step6;
    Else
        GoTo Step5
Step:6 If(I'm sheduled)
    GoTo Step7;
    Else
        Count++;
        If(Count=> threshold_count)
            GoTo Step1;
        Else
            GoTo Step5;
}
Step:7 I'm joining the network
Finish;

```

3. SIMULATION AND TEST

In this study, Riverbed (Opnet) [12] Modeler simulation environment has been preferred because it has many advantages such as an advanced graphical interface, hierarchical modeling, and simultaneously simulations with multiple inputs to test the algorithm. It also supports from very small networks to very large networks [13].



Figure 3. Scenario of 200 node in Riverbed Modeler network project environment (3200mX3200m)

In this study, the node counts have been defined as 25, 50, 100, 200, 400, 600, 800, 1000 to test the effect of node counts on network joining. The maximum coverage area of both CNs and SNs is set to 120m. Both CNs and nodes are randomly distributed in the simulation area. Thus, a multi-hop structure has been achieved. There are seven CNs in the network and each uses a different schedule channel. In order for the algorithm to have a flexible structure, the SNs configured to listen to all the CN channels. Network environments have been shown in Figure 3.

4. CONCLUSION

The proposed algorithm has shown that all nodes in all scenarios join to the network. As the number of nodes in the scenario increases, hop count increases, so the network joining latency have increased. In Figure 4, it has been shown that status of the node joining to the network for SN side in all scenarios. In Figure 5, it has been shown that status of the node joining to the network for CN side in all scenarios. As seen from the figures the differences between graphs are very small. As the number of nodes increases, the fluctuations in joining network have increased. Even if there is a delay as seen in the figures, all the nodes have joined the network.

Our main goal is not to join the nodes to the CN neither early nor late, but all nodes connect directly or indirectly (relay) to the CN.

Another goal of developed algorithm in this study is to ensure the continuity of connected nodes. It is obvious that the fast reconnection of nodes that are disconnected from the network is very important for the performance of the network. Looking at the results, it is seen that the proposed algorithm produces realistic results close to one hundred percent in network joining.

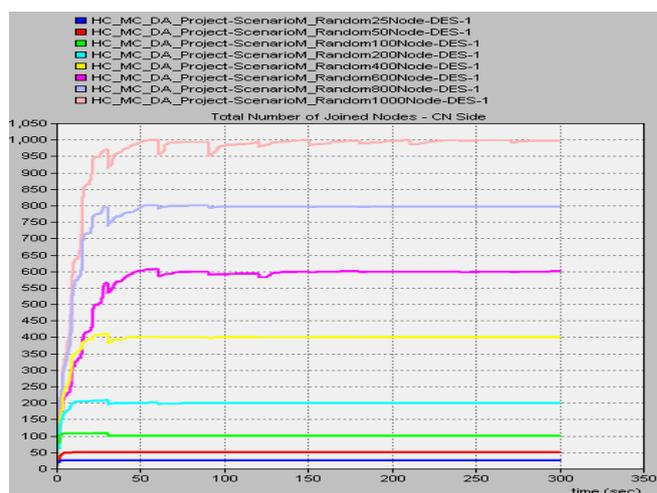


Figure 4. Total Number of Joined Nodes - ND Side

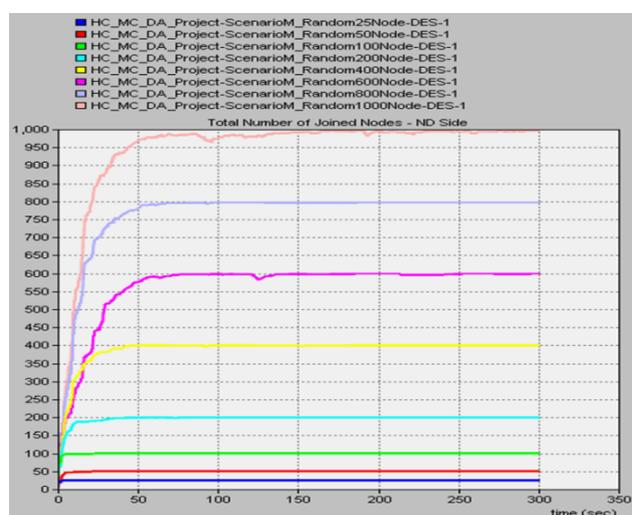


Figure 5. Total Number of Joined Nodes - CN Side

In this study, network joining algorithm considered for the non-mobile WSN networks and performance analysis has performed in the simulation environment. In the future, it is planned to adapt this algorithm for WSN networks consisting of mobile sensor nodes.

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BIOGRAPHIES

Musa Çıbuk obtained his bachelors' degree in the Electrical-Electronic Engineering Department at Firat University, Elazığ/Turkey in 1997, his master degree in the Electrical-Electronic Engineering Department at Firat University, Elazığ/Turkey in 2002 and PhD degree in the Electrical-Electronic Engineering Department at Firat University, Elazığ/Turkey in 2009. He is still an academic member of the Computer Engineering Department at Bitlis Eren University. His major areas of interests are: Wireless Communication, Internet, Server Systems, Sensor Networks and Image Processing.

Davut Arı obtained his BSc degree in computer engineering from Harran University (HRU) in 2014. He received the MSc diploma in Electric Engineering from the Bitlis Eren University in 2018. Her research interests are Computational Intelligence, Genetic Programming, Wireless Sensor Networks. In 2015 he joined the Computer Engineering, Faculty of Engineering And Architecture, Bitlis Eren University as a research assistant.

Fikri Ağgün obtained his bachelors' degree in the Computer Engineering Department at Selçuk University, Konya/Turkey in 2001, his master degree in the Biometry and Genetic Department at Yüzüncü Yıl University, Van/Turkey in 2011, Phd Study is resuming. He is still an academic member of the Informatics Department at Bitlis Eren University. His major areas of interests are: Computer Networks, Wireless Sensor Networks and Vehicular Adhoc Networks.